Updated S 09/360,542

L Number	Hits	Search Text	DB	Time stamp
1	15753	application adj program$4 adj interface	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:44
2	30774	automatic$4 with (backup or back-up or recover$6 or restor$6 or (back adj1 up))	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:37
3	23053	execut$6 near2 state$	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:45
4	0	(application adj program$4 adj interface) same (automatic$4 with (backup or back-up or recover$6 or restor$6 or (back adj1 up))) same (execut$6 near2 state$)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:28
5	97	(application adj program$4 adj interface) and (automatic$4 with (backup or back-up or recover$6 or restor$6 or (back adj1 up))) and (execut$6 near2 state$)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:28
6	10	(application adj program$4 adj interface) same (automatic$4 with (backup or back-up or recover$6 or restor$6 or (back adj1 up)))	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:29
7	0	((application adj program$4 adj interface) same (automatic$4 with (backup or back-up or recover$6 or restor$6 or (back adj1 up)))) and (execut$6 near2 state$)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:29
8	39	(application adj program$4 adj interface) same (execut$6 near2 state$)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:29
9	0	((application adj program$4 adj interface) same (execut$6 near2 state$)) and (automatic$4 with (backup or back-up or recover$6 or restor$6 or (back adj1 up)))	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:29
10	14	((application adj program$4 adj interface) same (execut$6 near2 state$)) and (backup or back-up or recover$6 or restor$6 or (back adj1 up))	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:31

11	34069	(application adj program$4 adj interface) or API	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:44
12	13338	execut$6 near1 state$	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:45
13	14992	(backup or back-up or recover$6 or restor$6 or (back adj1 up)) near2 (program$ or software or application)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:47
14	70	((application adj program$4 adj interface) or API) and (execut$6 near1 state$) and ((backup or back-up or recover$6 or restor$6 or (back adj1 up)) near2 (program$ or software or application))	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:48
15	0	((application adj program$4 adj interface) or API) same (execut$6 near1 state$) same ((backup or back-up or recover$6 or restor$6 or (back adj1 up)) near2 (program$ or software or application))	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:48
16	17008	((application adj program$4 adj interface) or API) same e ((backup or back-up or recover$6 or restor$6 or (back adj1 up)) near2 (program$ or software or application))	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:49
17	123	((application adj program$4 adj interface) or API) same ((backup or back-up or recover$6 or restor$6 or (back adj1 up)) near2 (program$ or software or application))	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:49
18	1	(((application adj program$4 adj interface) or API) same ((backup or back-up or recover$6 or restor$6 or (back adj1 up)) near2 (program$ or software or application))) and (execut$6 near1 state$)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:49
19	3569	automatic$4 adj2 recover$4	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:51
20	124	(automatic$4 adj2 recover$4) same restor$6	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 11:52

21	1	((automatic$4 adj2 recover$4) same restor$6) and ((application adj program$4 adj interface) or API)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:14
22	136	(automatic$4 adj2 recover$4) and (application adj program$4 adj interface)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:14
23	0	((automatic$4 adj2 recover$4) same restor$6) and (application adj program$4 adj interface)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:15
24	78	(automatic$4 adj2 recover$4) and restor$6 and (application adj program$4 adj interface)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:19
25	18	((automatic$4 adj2 recover$4) and restor$6 and (application adj program$4 adj interface)) and (714/$).ccls.	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:16
26	32	(automatic$4 adj2 recover$4) and (restor$6 adj2 (program or software or application))	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:40
27	4	((automatic$4 adj2 recover$4) and (restor$6 adj2 (program or software or application))) and (application adj program$4 adj interface)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:21
28	3734	(backup or back-up or (back$4 adj1 up)) near2 (software$ or program$ or application$)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:57
29	196	((backup or back-up or (back$4 adj1 up)) near2 (software$ or program$ or application$)) and (restor$6 adj2 (program or software or application))	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:40
31	9526	(hardware or execut$6) adj1 state$	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:41

32	1	((((backup or back-up or (back$4 adj1 up)) near2 (software$ or program$ or application$)) and (restor$6 adj2 (program or software or application))) and (application adj program$4 adj interface)) and ((hardware or execut$6) adj1 state$)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:42
30	16	(((backup or back-up or (back$4 adj1 up)) near2 (software$ or program$ or application$)) and (restor$6 adj2 (program or software or application))) and (application adj program$4 adj interface)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:56
33	1		USPAT	2004/03/11 12:54
34	1		USPAT	2004/03/11 12:54
35	1		USPAT	2004/03/11 12:55
36	1		USPAT	2004/03/11 12:55
37	1		USPAT	2004/03/11 12:55
38	1		USPAT	2004/03/11 12:55
39	20539	restor$6.ti.	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:58
40	113580	recover$6.ti.	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:58
41	17718	(backup or back-up or (back$4 adj1 up)).ti.	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:58
42	0	restor$6.ti. and recover$6.ti. and ((backup or back-up or (back$4 adj1 up)).ti.) and (application adj program$4 adj interface)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:58
43	112780	restor$6.ab.	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:58
44	327436	recover$6.ab.	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:58

45	49561	(backup or back-up or (back$4 adj1 up)).ab.	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:59
46	53855	(application adj program$4 adj interface) and restor$6.ab. and recover$6.ab. adn ((backup or back-up or (back$4 adj1 up)).ab.)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 12:59
47	4	(application adj program$4 adj interface) and restor$6.ab. and recover$6.ab. and ((backup or back-up or (back$4 adj1 up)).ab.)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 13:01
48	2	(registry or catalog) and ((application adj program$4 adj interface) and restor$6.ab. and recover$6.ab. and ((backup or back-up or (back$4 adj1 up)).ab.))	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 13:03
49	30	(registry or catalog) and recover$6.ab. and ((backup or back-up or (back$4 adj1 up)).ab.)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 13:03
50	4	((registry or catalog) and recover$6.ab. and ((backup or back-up or (back$4 adj1 up)).ab.)) and ((application adj program$4 adj interface) or API)	USPAT; US-PGPUB; EPO; JPO; DERWENT; IBM_TDB	2004/03/11 13:04

US-PAT-NO: 6564215

DOCUMENT-IDENTIFIER: US 6564215 B1

TITLE: Update support in database content management

---------- KWIC ---------

Abstract Text - ABTX (1):

A computer system updates a data object that is maintained in data storage external to a database management system (DBMS), after receiving an update request from a DBMS client for the data object, by first scheduling the update request with the DBMS to provide access to the external data object, then initiating a subtransaction in the DBMS for the update request, updating the data object with an in-place update action at the external data storage to thereby produce an updated data object and also updating the DBMS metadata of the data object, next appending information relating to type and time of the update action in an object version table, and then executing a **backup** operation of the updated data object. This permits update-in-place operations on the external data object, under supervision of the DBMS. The system thereby supports update-in-place operations on external data with access control, **backup and recovery,** and transaction consistency in accordance with a database management system, while avoiding large copy operations that would consume network resources.

Detailed Description Text - DETX (2):

FIG. 1 is a representation of the functional components contained in a computer system 100 constructed in accordance with the present invention. The computer system includes the datalinks engine 102 of a database management system (DBMS) 104 at a first node of a computer network, and two file sites 106, 108 communicating with the DBMS over the network 110. Data objects 109 comprising files are stored at the file sites 106, 108 and are considered external to the DBMS 104. A DBMS client 112 at another network node

113 also
communicates with the DBMS 104 over the network. In accordance with the
invention, a computer program application 114 at the client node 113 updates a
data object 109 stored at one of the file sites 106, 108 by directly accessing
and updating the data object in the external store through a file system
application program interface (API) 116 at each respective file site without
first making a copy of the data object. The client application 114 also
accesses metadata of the data object through an SQL interface provided by the
DBMS 104 to maintain consistency between the data object and its metadata.
This technique permits update-in-place operations on the data object, under
supervision of the DBMS 104. In this way, the computer system 100 supports
update operations on external data with access control, backup and recovery,
and transaction consistency in accordance with the DBMS, while avoiding large
copy operations that would consume network resources and also maintaining the
same object name from the file system perspective.

Detailed Description Text - DETX (5):
In a conventional DBMS system, the file will be checked out via the DBMS
104, and a copy of a checked out data object 109 will be sent from the file
site 106 to the client 112, where the application 114 can be used for update
operations on the copy. Following the update operations, the client 112
returns (checks in) the updated data object copy to the file site 106,
whereupon the updated data object 109 is stored. As noted above, however, the
computer system 100 of the present invention can support update-in-place
operations by providing a system having a DBMS in which client applications 114
can directly update data objects through the **API** 116 of the file system at an
external file site 106. The applications 114 can be any one of a variety of
client applications, including applications that provide word processing,
spreadsheet, database, and Internet-protocol web browser functions. The direct
updating from the client application to the external file system **API** is
represented in FIG. 1 by the dashed line 121.

Detailed Description Text - DETX (24):
After the DLFM processing of the Update_Pending request, the DLFS continues the update-in-place processing by passing the file open request to the native file system at the external store. The requested file is then opened by the native file system and a file pointer is returned to the client. This operation is represented by the FIG. 4 flow diagram box numbered 410. All subsequent update processing on the requested file, including write and read operations, will then be carried out by the client application without any further intervention by either the DLFM or DLFS. Such subsequent update processing can be carried out using the **API** of the external store file system. These update operations are represented in FIG. 4 by the flow diagram box numbered 412.

Detailed Description Text - DETX (33):
When a "linked" file (a file whose access is under database system control) is updated, the preferred embodiment makes an archive copy of the updated file. Each archived version of the updated file is associated with a unique identifier and the information is stored in the File_Version table. In a restore operation, after restoring the database to a condition at a specific point in time in the past, the corresponding files would also have to be restored from the archive to match the restored database state, or condition. To do so, the database system can start a reconciliation process to synchronize the database state and the file system state. The following are the steps the system performs for such reconciliation processing. (a) A database agent sets an identifier (called db_state_id) for the current database state which can be implemented by a time stamp or tail LSN. (b) The database agent scans the **catalog** tables to identify all user tables having one or more columns with a Datalink data type that support the update-in-place processing described herein. (c) For each table found, the database agent scans the table and extracts the datalink column value (URL) from each record. (d) The agent inserts the extracted URL into a message block and sends it, along with the database state identifier db_state_id, to a DLFM daemon process. (e)

The DLFM daemon process uses the URL and db_state_id to look up the corresponding entry in a DLFM table (the File_Version table) and then extracts the archive file name from the DLFM table entry. When there is more than one entry with the same URL, the one with the greatest value of unique recovery id that is smaller than or equal to the received db_state_id is selected. (f) The DLFM daemon process then restores the file to the file system from the archive device if the file is not already in the file system.

Detailed Description Text - DETX (35):

FIG. 7 is a flow diagram that illustrates the processing steps executed by the computer processing system 100 of FIG. 1 to implement the coordinated database restore operation. The flow diagram illustrates the operation steps described above. Thus, in the first processing step, a database agent sets an identifier for the current database state, which can include time stamp data. This processing is represented by the flow diagram box numbered 702, and corresponds to step (a) above. Next, the database agent scans the **catalog** tables to identify all user tables having one or more columns with Datalink data types that support the update-in-place processing described herein. This step is represented by box 704 in FIG. 7, and step (b) above. Next, for each table found, the database agent scans the table and extracts the URL from each record, as represented by box 706 and step (c). In the processing represented by box 708 of FIG. 7 (and step (d) above), the database agent inserts the extracted URL into a message block and sends it, along with the database state identifier db_state_id, to a DLFM daemon process.

US-PAT-NO: 6014437

DOCUMENT-IDENTIFIER: US 6014437 A

TITLE: Multi service platform architecture for telephone networks

---------- KWIC ---------

Detailed Description Text - DETX (41):
The call processors (CPs)--consisting of a pair of RISC System/6000 computers based on the IBM RISC System/6000 workstation--execute layered software (see discussion of software architecture in next section below) that includes the logic determining how calls are treated and processed. Both CPs are connected to and share the same physical disks for data storage, and data stored on the disks is available to both. The CPs are configured in a high availability arrangement wherein both are active and processing calls, but which can fall back to continued operation with a single CP when necessary. The CPs use "mirrored" disks to ensure reliability and provide quick access to data. Each processor has a CD-ROM drive and tape drive for loading its respective operating system **software and other software and for backing up** and restoring information accessed principally through the mirrored disks.

Detailed Description Text - DETX (47):
FIG. 3 provides an overview of the software architecture presently contemplated for the MSP/6000 system. The software is architected to operate in three layers--a service application management layer 30, a service application layer 31, and a communications and management layer 32--based on a platform 33 whose major component is the AIX/6000 Operating System commonly used in RISC/6000 processors. An **application programming interface** associated with service creation (Service Creation API) 35 links layers 31 and 32.

Detailed Description Text - DETX (48):
The Application Management layer 30--which is intended to be shared by multiple MSP/6000 nodes, and may be located physically and geographically separate from any and all of the nodes--consists of an element

management
system (EMS) 40 and a service creation environment (SCE) 41. SCE 41 provides
both a graphical user interface (GUI) and a text-based application oriented
language (AOL) suited for the development, debugging and testing of telephony
service application logic (e.g. a language based on the DirectTalk/6000 state
table service language). Using such language, developers can create
applications that play back selected pre-recorded announcements serving as
prompts to a user; such prompts being either prerecorded and edited from a tape
cassette unit or recorded live using a microphone. EMS 40 includes tools for
distributing, backing up and **restoring service applications** (including recorded
prompts). The recorded prompts are distributed/restored to call processors in
(multiple) MSP/6000 systems and transferred by the latter processors to voice
peripherals in respective systems. Tools for backing up and restoring the
prompts are distributed to MOC processors in respective systems and accessed
therefrom when needed.